DTU

Verifying a Sequent Calculus Prover

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Introduction

- A sound and complete prover for first-order logic with functions
- Based on a sequent calculus
- All proofs are formally verified in Isabelle/HOL
- Human-readable proof certificates

Background

- Formalized metatheory for non-trivial sequent calculus provers
- Formal verification of an executable prover
- Novel analytic proof technique for completeness
- Verifiable and human-readable proof certificates
- A prover for the SeCaV system

Sample SeCaV Proof Rules

$$\frac{\text{Neg } p \in z}{\Vdash p, z} \ \textsc{Basic} \qquad \frac{\Vdash z \qquad z \subseteq y}{\Vdash y} \ \textsc{Ext} \qquad \frac{\Vdash p, z}{\Vdash \text{Neg (Neg } p), z} \ \textsc{NegNeg}$$

$$\frac{\Vdash p, q, z}{\Vdash \text{Dis } p\ q, z} \ \textsc{AlphaDis} \qquad \frac{\Vdash \text{Neg } p, z \qquad \Vdash \text{Neg } q, z}{\Vdash \text{Neg (Dis } p\ q), z} \ \textsc{BetaDis}$$

$$\frac{\Vdash p\,[\text{Var } 0/t], z}{\Vdash \text{Exi } p, z} \ \textsc{GammaExi}$$

$$\frac{\Vdash \text{Neg } (p\,[\text{Var } 0/\text{Fun } i\ []]), z \qquad i \text{ fresh}}{\Vdash \text{Neg (Exi } p), z} \ \textsc{DeltaExi}$$

Prover I

- SeCaV rules affect *one formula* at a time
- Our prover rules affect *every applicable formula* at once
- We copy Gamma formulas and remember all terms on the branch
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- Rules affect disjoint formulas
- So we can apply them in any order
- We apply rules *fairly* and repeatedly
- So we never miss out on a proof

Prover II

- We rely on the abstract completeness framework by Blanchette, Popescu and Traytel
- We need to fix a *stream* of rules from the beginning
- Proof attempts are *coinductive trees* grown by applying these rules
- If a tree cannot be grown further, we found a proof
- A function gives the *child sequents* representing the subgoals left after applying a rule
- We export code to Haskell to obtain an executable prover

Prover — proof example

$$\dfrac{}{\begin{array}{c}\text{Neg (Uni (Con } P(0)\ Q(0))), \text{Neg } P(0), \text{Neg } Q(0),\\ \text{Neg } P(a), \text{Neg } Q(a), P(a)\end{array}}\ \textsc{Basic}$$

$$\dfrac{}{\begin{array}{c}\text{Neg (Uni (Con } P(0)\ Q(0))), \text{Neg (Con } P(0)\ Q(0)),\\ \text{Neg (Con } P(a)\ Q(a)), P(a)\end{array}}\ \textsc{AlphaCon}$$

$$\dfrac{}{\begin{array}{c}\text{Neg (Uni (Con } P(0)\ Q(0))), \text{Neg (Con } P(0)\ Q(0)),\\ \text{Neg (Con } P(a)\ Q(a)), P(a)\end{array}}\ (\alpha)$$

$$\dfrac{}{\text{Neg (Uni (Con } P(0)\ Q(0))), P(a)}\ \textsc{GammaUni}$$

$$\dfrac{}{\text{Neg (Uni (Con } P(0)\ Q(0))), P(a)}\ (\alpha, \delta, \beta)$$

$$\dfrac{}{\text{Imp (Uni (Con } P(0)\ Q(0)))\ P(a)}\ \textsc{AlphaImp}$$

$$\dfrac{}{\text{Imp (Uni (Con } P(0)\ Q(0)))\ P(a)}\ (\textsc{NegNeg})$$

Prover — certificate example

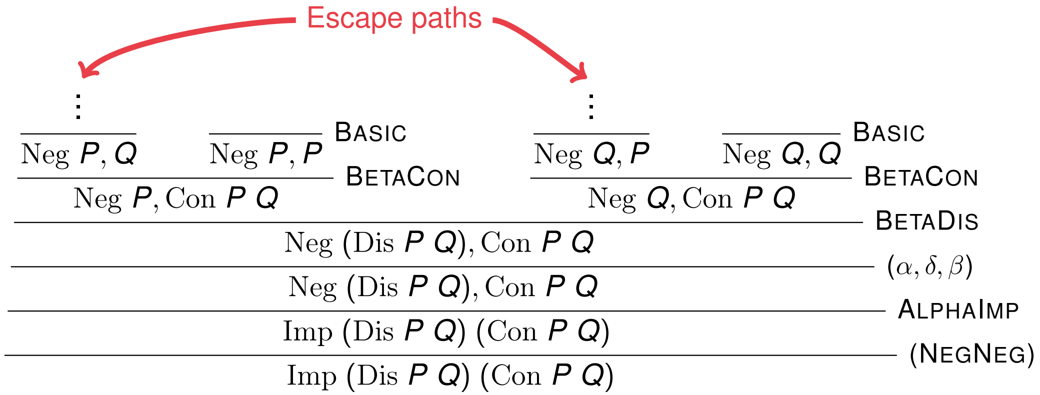
```
Imp (Uni (Con (P [0]) (Q [0]))) (P [a])

AlphaImp
  Neg (Uni (Con (P [0]) (Q [0])))
  P [a]
Ext
  Neg (Uni (Con (P [0]) (Q [0])))
  Neg (Uni (Con (P [0]) (Q [0])))
  P [a]
GammaUni [0]
  Neg (Con (P [0]) (Q [0]))
  Neg (Uni (Con (P [0]) (Q [0])))
  P [a]
```

```
Ext
  Neg (Uni (Con (P [0]) (Q [0])))
  Neg (Uni (Con (P [0]) (Q [0])))
  P [a]
  Neg (Con (P [0]) (Q [0]))
GammaUni [a]
  Neg (Con (P [a]) (Q [a]))
  Neg (Uni (Con (P [0]) (Q [0])))
  P [a]
  Neg (Con (P [0]) (Q [0]))
Ext
  Neg (Con (P [0]) (Q [0]))
  Neg (Con (P [a]) (Q [a]))
  P [a]
  Neg (Uni (Con (P [0]) (Q [0])))
AlphaCon
  Neg (P [0])
  Neg (Q [0])
  Neg (Con (P [a]) (Q [a]))
  P [a]
  Neg (Uni (Con (P [0]) (Q [0])))
```

```
Ext
  Neg (Con (P [a]) (Q [a]))
  P [a]
  Neg (Uni (Con (P [0]) (Q [0])))
  Neg (P [0])
  Neg (Q [0])
AlphaCon
  Neg (P [a])
  Neg (Q [a])
  P [a]
  Neg (Uni (Con (P [0]) (Q [0])))
  Neg (P [0])
  Neg (Q [0])
Ext
  P [a]
  Neg (Uni (Con (P [0]) (Q [0])))
  Neg (P [0])
  Neg (Q [0])
  Neg (P [a])
  Neg (Q [a])
Basic
```

Prover — escape path example



Soundness I

- If our prover returns a proof, we can build a SeCaV proof
- The SeCaV proof system is sound, so the prover is sound
- We use the abstract soundness framework by Blanchette et al.
- If the *children* of a sequent all have SeCaV proofs, so does the sequent

Soundness II

If the *children* of a sequent all have SeCaV proofs, so does the sequent:

1. Assume all child sequents have a proof
2. Induction on sequent: use appropriate SeCaV rule for each formula

Example: Our sequent looks like $\mathrm{Dis}\ P\ Q, \ldots$, so $P, Q, \ldots$ is a child sequent with a SeCaV proof. We apply the AlphaDis rule to prove the sequent using the proof of $\Vdash P, Q, \ldots$ (and possibly some reordering).

$$\dfrac{\dfrac{\ldots \qquad \dfrac{\vdots}{\Vdash P, Q, \ldots}\ \textsc{Assumption} \qquad \ldots}{\Vdash \mathrm{Dis}\ P\ Q, \ldots}\ \textsc{AlphaDis}}{\vdots}$$

Completeness

- Framework: prover either produces a finite, well formed proof tree or an infinite tree with a saturated escape path
- Need to show that root sequent of a saturated escape path is not valid:
  - Formulas on saturated escape paths form Hintikka sets
  - Hintikka sets induce a well formed countermodel
- . . . so valid sequents result in finite, well formed proof trees

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- ...so valid sequents result in finite, well formed proof trees

HERE BE DRAGONS

(need to build a bounded countermodel over only the terms in the sequent and ensure functions stay inside its domain)

Bounded semantics

- In a completeness proof for a *calculus* we can assume that Gamma formulas are instantiated with all possible terms
- Thus, we can build a countermodel in the full Herbrand domain
- Our prover only uses terms from the given sequent (and fresh ones)
- So we must build a *bounded* countermodel over this restricted domain
- We must ensure that our function denotation stays inside this domain

Subtypes fail us

- The SeCaV semantics represents the domain as a type variable.
- We cannot build the subtype of terms from a *local* sequent (yet?[1])
- So we represent the domain as an explicit parameter to the semantics
- We have $u, E, F, G \models \mathrm{Uni}\ P$ iff $u, E, F, G \models P(x)$ for all $x \in u$
- We reprove soundness of SeCaV under this (u)semantics

[1] Kunčar and Popescu ITP 2014

Hintikka sets

- We always need at least one term

$$\mathit{terms}\ H \equiv \mathit{if}\ (\bigcup p \in H.\ \mathit{set}\ (\mathit{subtermFm}\ p)) = \{\}\ \mathit{then}\ \{\mathit{Fun}\ 0\ []\}$$
$$\mathit{else}\ (\bigcup p \in H.\ \mathit{set}\ (\mathit{subtermFm}\ p))$$

- To quantify over in our Hintikka sets

locale *Hintikka* =
  fixes $H :: \mathit{fm\ set}$
  assumes
    Basic: $\mathit{Pre}\ n\ \mathit{ts} \in H \Longrightarrow \mathit{Neg}\ (\mathit{Pre}\ n\ \mathit{ts}) \notin H$ **and**
    $\mathit{AlphaDis}$: $\mathit{Dis}\ p\ q \in H \Longrightarrow p \in H \wedge q \in H$ **and**
    $\mathit{BetaDis}$: $\mathit{Neg}\ (\mathit{Dis}\ p\ q) \in H \Longrightarrow \mathit{Neg}\ p \in H \vee \mathit{Neg}\ q \in H$ **and**
    $\mathit{GammaExi}$: $\mathit{Exi}\ p \in H \Longrightarrow \forall t \in \mathit{terms}\ H.\ \mathit{sub}\ 0\ t\ p \in H$ **and**
    $\mathit{DeltaExi}$: $\mathit{Neg}\ (\mathit{Exi}\ p) \in H \Longrightarrow \exists t \in \mathit{terms}\ H.\ \mathit{Neg}\ (\mathit{sub}\ 0\ t\ p) \in H$ **and**
    $\vdots$

Bounded countermodel

- We carefully build the countermodel

  $E\ S\ n \equiv \mathit{if}\ \mathit{Var}\ n \in \mathit{terms}\ S\ \mathit{then}\ \mathit{Var}\ n\ \mathit{else}\ \mathit{SOME}\ t.\ t \in \mathit{terms}\ S$
  $F\ S\ i\ l \equiv \mathit{if}\ \mathit{Fun}\ i\ l \in \mathit{terms}\ S\ \mathit{then}\ \mathit{Fun}\ i\ l\ \mathit{else}\ \mathit{SOME}\ t.\ t \in \mathit{terms}\ S$
  $G\ S\ n\ ts \equiv \mathit{Neg}\ (\mathit{Pre}\ n\ ts) \in S$
  $M\ S \equiv \mathit{usemantics}\ (\mathit{terms}\ S)\ (E\ S)\ (F\ S)\ (G\ S)$
- *terms* is downwards closed, so members evaluate to themselves

  $t \in \mathit{terms}\ S \Longrightarrow \mathit{semantics\text{-}term}\ (E\ S)\ (F\ S)\ t = t$
- We have a countermodel to any formula in a Hintikka set

  $\mathit{Hintikka}\ S \Longrightarrow (p \in S \longrightarrow \neg\ M\ S\ p) \wedge (\mathit{Neg}\ p \in S \longrightarrow M\ S\ p)$

Saturated escape paths form Hintikka sets

- Final step is to inspect the saturated escape paths
- We need to show that the formulas constitute a Hintikka set
- On paper, this follows straightforwardly from our rules
- In practice, it requires fiddly reasoning about the coinductive paths
- In the end: any saturated escape path has a (bounded) countermodel, contradicting the validity of its root sequent

Results and future work

- We have verified soundness and completeness in Isabelle/HOL
- Verification helped find actual bugs in our implementation
- The performance is limited, but optimizations are possible
- Generation of proof certificates is not (yet) verified
- Potentially consider extensions to the logic such as equality